

Week 11 - Monday

**COMP 2230**

# Last time

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- Trees
- Rooted trees
- Spanning trees

Questions?

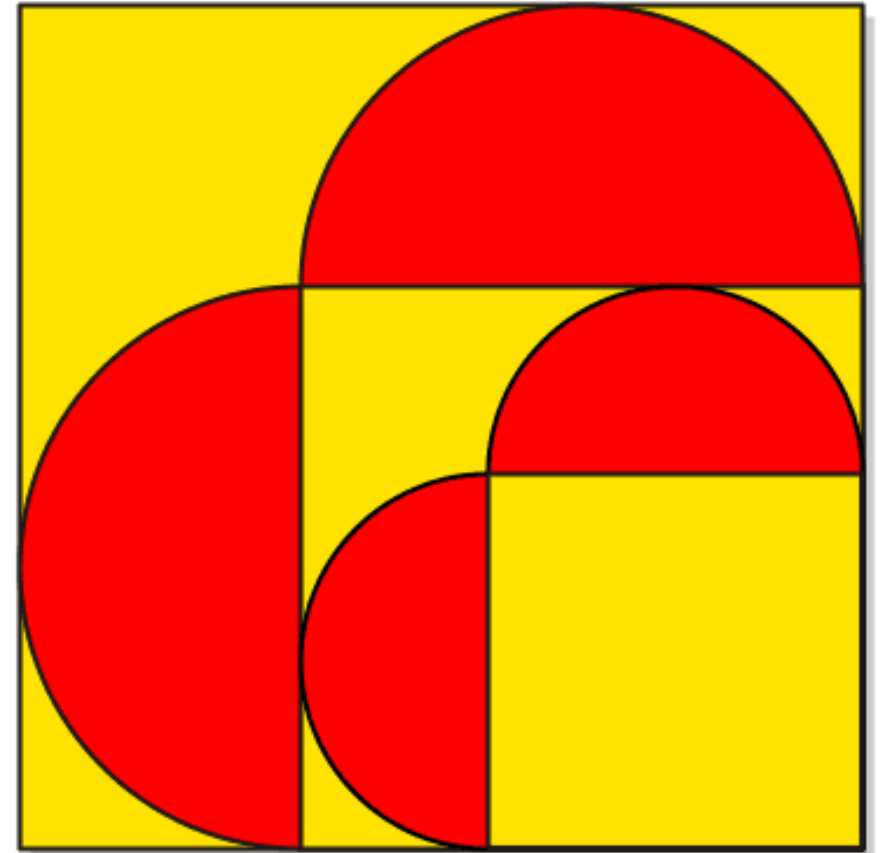
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# Assignment 5

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# Logical warmup

- Which has a bigger total area?
  - Yellow
  - Red



# Spanning Trees

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# Finding a minimum spanning tree

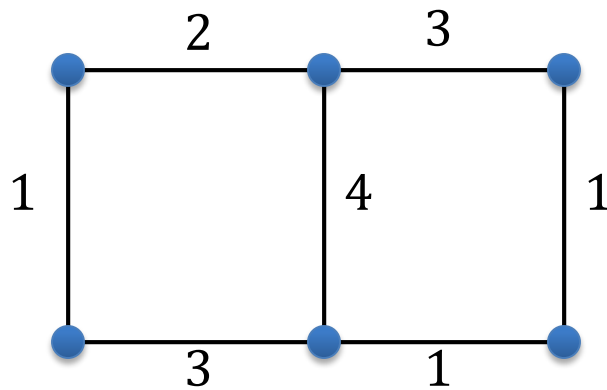
- Kruskal's algorithm gives an easy to follow technique for finding an MST on a weighted, connected graph
- Informally, go through the edges, adding the smallest one, unless it forms a circuit
- **Algorithm:**
  - Input: Graph  $G$  with  $n$  vertices
  - Create a subgraph  $T$  with all the vertices of  $G$  (but no edges)
  - Let  $E$  be the set of all edges in  $G$
  - Set  $m = 0$
  - While  $m < n - 1$ 
    - Find an edge  $e$  in  $E$  of least weight
    - Delete  $e$  from  $E$
    - If adding  $e$  to  $T$  doesn't make a circuit
      - Add  $e$  to  $T$
      - Set  $m = m + 1$
  - Output:  $T$

# Prim's algorithm

- Prim's algorithm gives another way to find an MST
- Informally, start at a vertex and add the next closest node not already in the MST
- **Algorithm:**
  - Input: Graph  $G$  with  $n$  vertices
  - Let subgraph  $T$  contain a single vertex  $v$  from  $G$
  - Let  $V$  be the set of all vertices in  $G$  except for  $v$
  - For  $i$  from 1 to  $n - 1$ 
    - Find an edge  $e$  in  $G$  such that:
      - $e$  connects  $T$  to one of the vertices in  $V$
      - $e$  has the lowest weight of all such edges
    - Let  $w$  be the endpoint of  $e$  in  $V$
    - Add  $e$  and  $w$  to  $T$
    - Delete  $w$  from  $V$
  - Output:  $T$

# Prim fights Kruskal

- Apply Kruskal's algorithm to the graph below
- Now, apply Prim's algorithm to the graph below
- Is there any other MST we could make?



# Graphing Functions and Big-O, Big-Omega, and Big-Theta Notations

Three-Sentence Summary

# Graphing Functions

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# Graphing functions

- I am assuming that you've been graphing functions since about 6<sup>th</sup> or 7<sup>th</sup> grade
- Still, the formal definition of a graph of a function is:
  - Let  $f$  be a real-valued function of a real variable
  - The **graph of  $f$**  is the set of all points  $(x, y)$  in the Cartesian coordinate plane such that  $x$  is in the domain of  $f$  and  $y = f(x)$

# Power functions

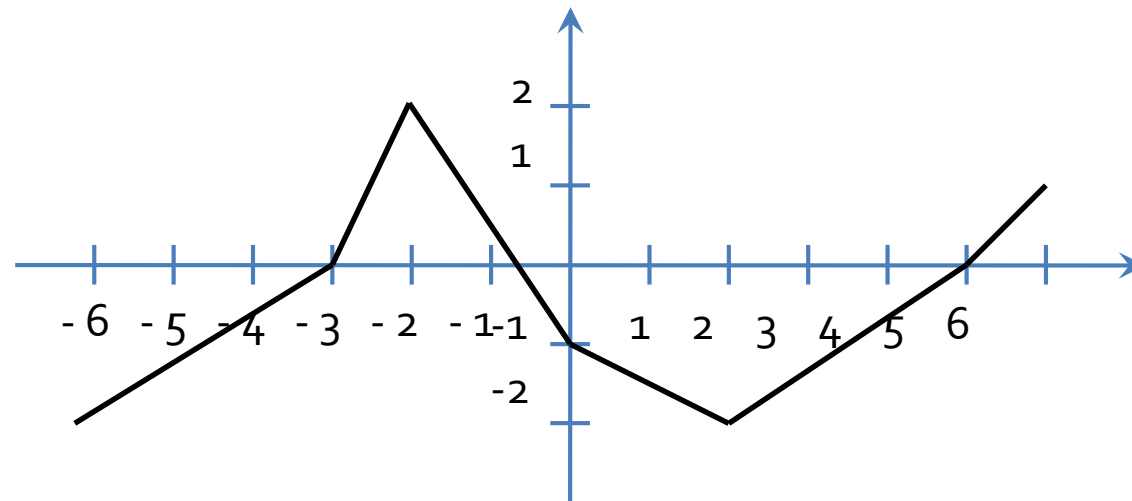
- In computer science, we are often interested in **power functions of  $x$**  written  $p_a(x)$  where  $p_a(x) = x^a$  where  $x$  and  $a$  are nonnegative
- Power functions are the building blocks of polynomial functions
- Graph the following on the same coordinate axes
  - $p_0$
  - $p_{\frac{1}{2}}$
  - $p_1$
  - $p_2$

# Discontinuous functions

- Recall the definition of the floor of  $x$ :
  - $\lfloor x \rfloor$  = the largest integer that is less than or equal to  $x$
- Graph  $f(x) = \lfloor x \rfloor$
- Defining functions on integers instead of real values affects their graphs a great deal
- Graph  $p_1(x) = x, x \in \mathbb{R}$
- Graph  $f(n) = n, n \in \mathbb{N}$

# Multiples of functions

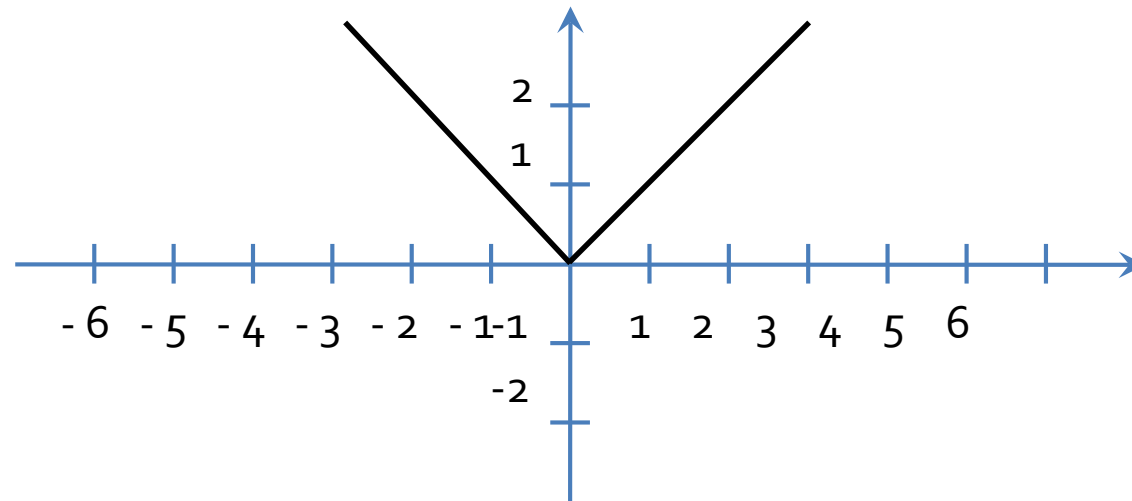
- There is a strong visual (and of course mathematical) correlation to a function that is the multiple of another function
- Examples:
  - $g(x) = x + 2$
  - $2 \cdot g(x) = 2x + 4$
- Given  $f$  graphed below, sketch  $2f$



# Absolute value

- Consider the absolute value function

- $f(x) = |x|$



- Left of the origin it is constantly **decreasing**
- Right of the origin it is constantly **increasing**

# Increasing and decreasing functions

- We say that  $f$  is **decreasing on the set**  $S$  iff for all real numbers  $x_1$  and  $x_2$  in  $S$ , if  $x_1 < x_2$ , then  $f(x_1) > f(x_2)$
- We say that  $f$  is **increasing on the set**  $S$  iff for all real numbers  $x_1$  and  $x_2$  in  $S$ , if  $x_1 < x_2$ , then  $f(x_1) < f(x_2)$
- We say that  $f$  is an **increasing** (or **decreasing**) function iff  $f$  is increasing (or decreasing) on its entire domain
- A positive multiple of an increasing function is increasing
- Virtually all running time functions are increasing functions

# Asymptotic Notations

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# Growth of functions

- Mathematicians worry about the growth of various functions
- They usually express such things in terms of limits, maybe with derivatives
- We are focused primarily on functions that bound running time spent and memory consumed
- We just need a rough guide
- We want to know the **order** of the growth

# Definitions

- Let  $f$  and  $g$  be real-valued functions defined on the same set of nonnegative real numbers
- $f$  is of order at least  $g$ , written  $f(x)$  is  $\Omega(g(x))$ , iff there is a positive  $A \in \mathbb{R}$  and a nonnegative  $a \in \mathbb{R}$  such that
  - $A|g(x)| \leq |f(x)|$  for all  $x > a$
- $f$  is of order at most  $g$ , written  $f(x)$  is  $O(g(x))$ , iff there is a positive  $B \in \mathbb{R}$  and a nonnegative  $b \in \mathbb{R}$  such that
  - $|f(x)| \leq B|g(x)|$  for all  $x > b$
- $f$  is of order  $g$ , written  $f(x)$  is  $\Theta(g(x))$ , iff there are positive  $A, B \in \mathbb{R}$  and a nonnegative  $k \in \mathbb{R}$  such that
  - $A|g(x)| \leq |f(x)| \leq B|g(x)|$  for all  $x > k$

# Using the notation

- Express the following statements using appropriate notation:
  - $10|x^6| \leq |17x^6 - 45x^3 + 2x + 8| \leq 30|x^6|$ , for  $x > 2$
  - $15|\sqrt{x}| \leq \left| \frac{15\sqrt{x}(2x+9)}{x+1} \right|$ ,  $x > 0$
  - $\left| \frac{15\sqrt{x}(2x+9)}{x+1} \right| \leq 15|\sqrt{x}|$ ,  $x > 6$
- Justify the following:
  - $\left| \frac{15\sqrt{x}(2x+9)}{x+1} \right|$  is  $\Theta(\sqrt{x})$

# Properties of $\Omega$ -, $O$ -, and $\Theta$ -notation

- Let  $f, g, h,$  and  $k$  be real-valued functions defined on the same set of nonnegative real numbers
  1.  $f(x)$  is  $\Omega(g(x))$  and  $f(x)$  is  $O(g(x))$  iff  $f(x)$  is  $\Theta(g(x))$
  2.  $f(x)$  is  $\Omega(g(x))$  iff  $g(x)$  is  $O(f(x))$
  3.  $f(x)$  is  $\Omega(f(x)), f(x)$  is  $O(f(x)),$  and  $f(x)$  is  $\Theta(f(x))$
  4. If  $f(x)$  is  $O(g(x))$  and  $g(x)$  is  $O(h(x))$  then  $f(x)$  is  $O(h(x))$
  5. If  $f(x)$  is  $O(g(x))$  and  $c$  is a positive real, then  $c \cdot f(x)$  is  $O(g(x))$
  6. If  $f(x)$  is  $O(h(x))$  and  $g(x)$  is  $O(k(x))$  then  $f(x) + g(x)$  is  $O(G(x))$  where  $G(x) = \max(|h(x)|, |k(x)|)$  for all  $x$
  7. If  $f(x)$  is  $O(h(x))$  and  $g(x)$  is  $O(k(x))$  then  $f(x) \cdot g(x)$  is  $O(h(x) \cdot k(x))$

# Ticket Out the Door

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# Upcoming

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# Next time...

- Analysis of algorithm efficiency I
- Exponential and logarithmic functions

# Reminders

- **Work on Assignment 5**
- Read 11.3 and 11.4
  - Prepare a three-sentence summary
  - Extra credit if you get called on